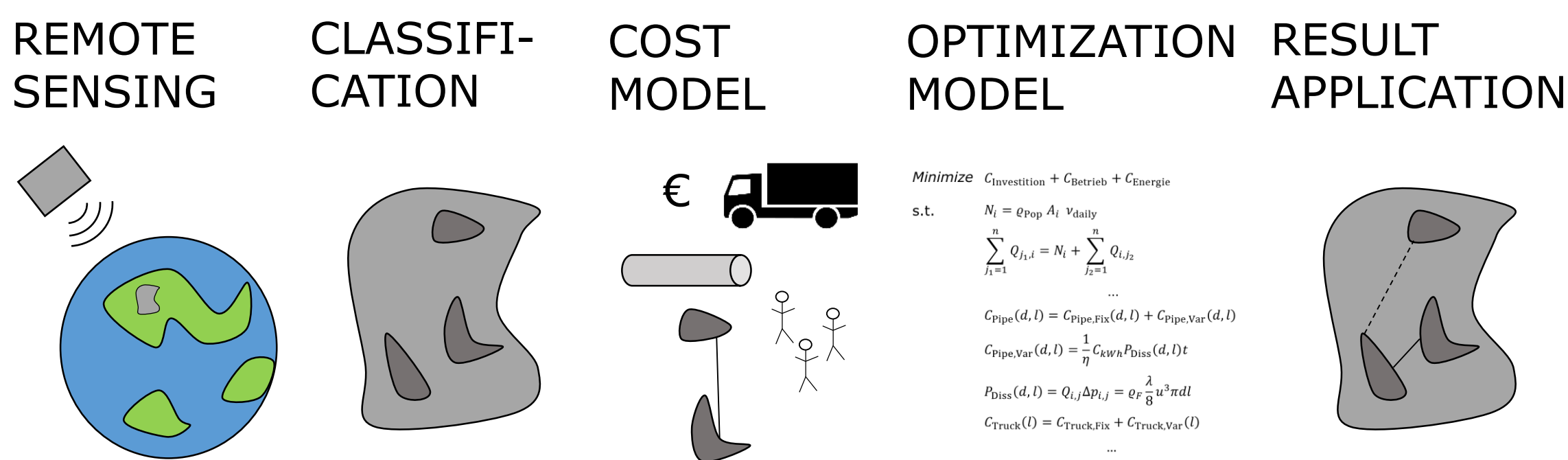


Optimizing water supply networks for slums in urban areas

Lea Rausch, Prof. Dr.-Ing. Peter F. Pelz

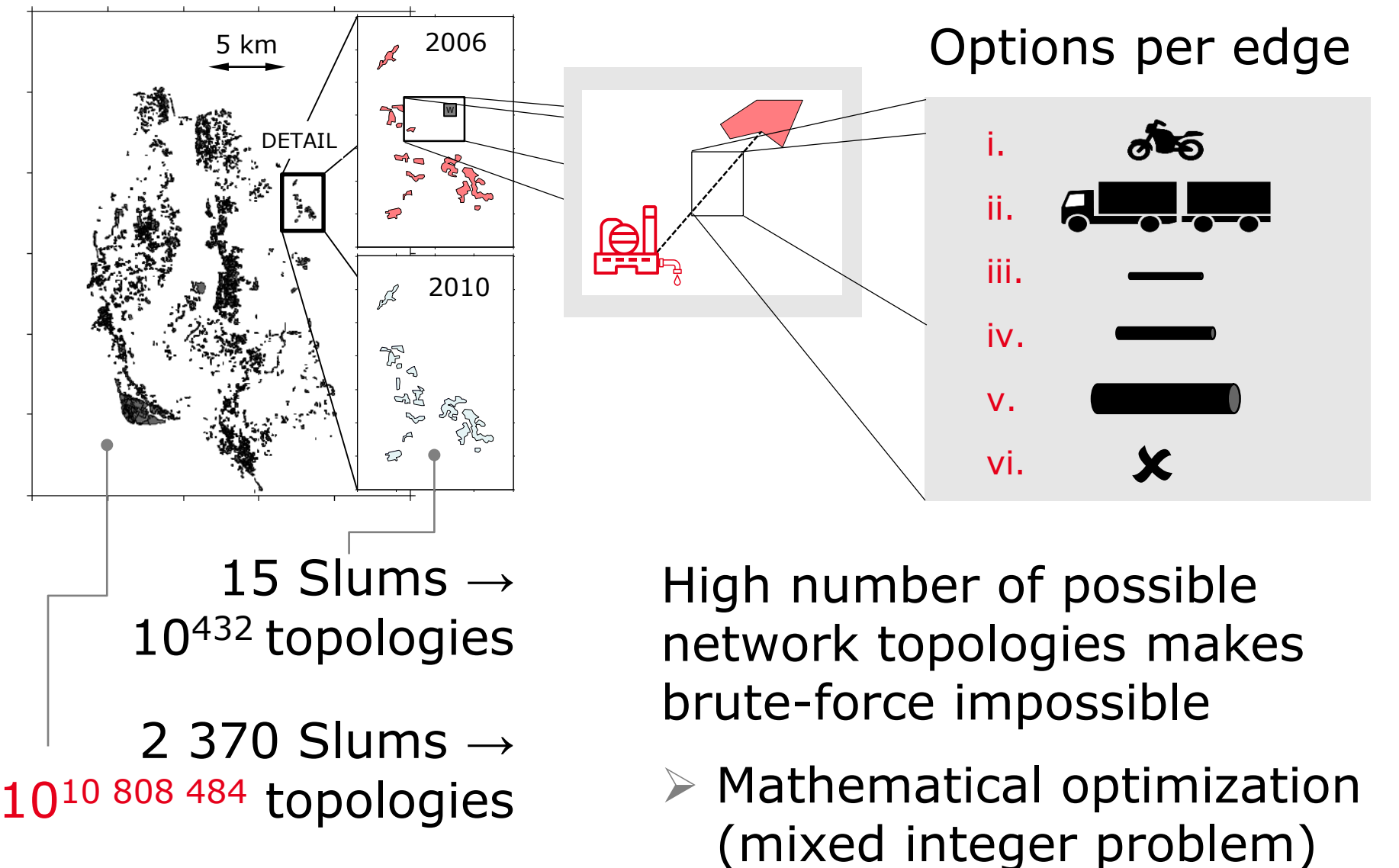
Developed Framework



Interdisciplinary approach to plan optimal water supply networks for slums in a large urban area

- Using remote sensing data to classify slums for demand estimate
- Combining different ways of transportation (e.g. pipes, vehicles)
- Minimizing total costs and CO₂-Emissions

Addressed Complication



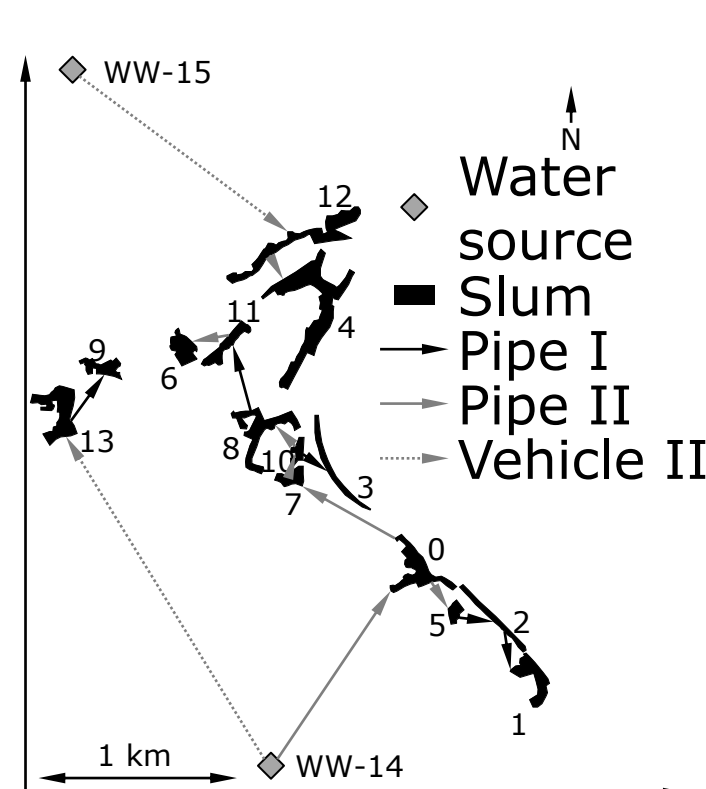
Generated Results

Optimization model formulated and solved as Mixed Integer Linear Problem (MILP). Further model extensions include

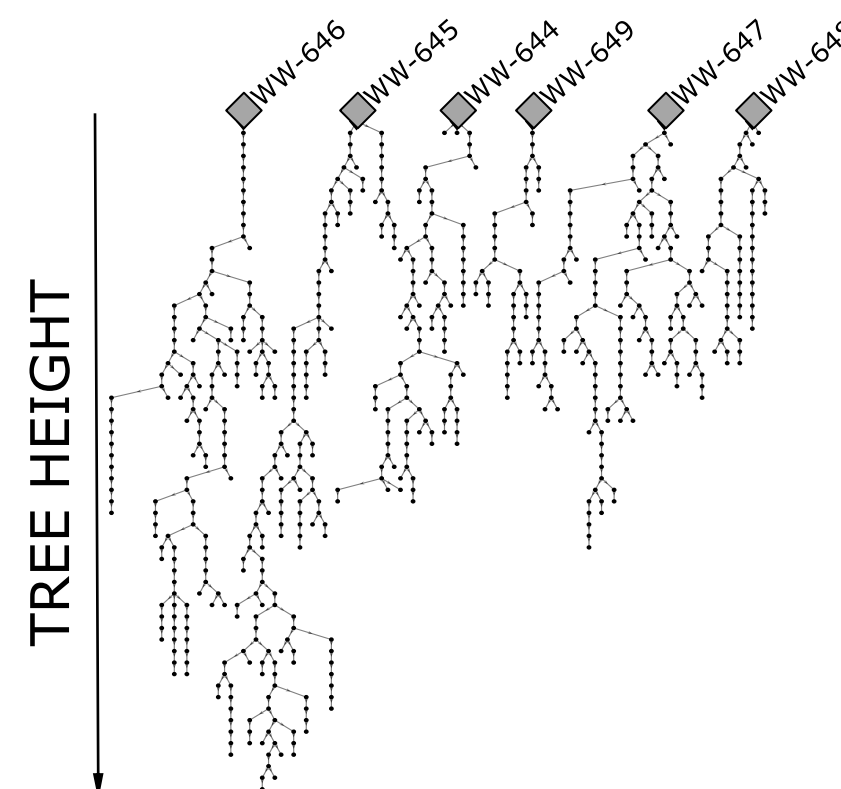
- Grey water network planning
- CO₂-Emission minimization

Developed solution methods address runtime issues with standard MILP solvers

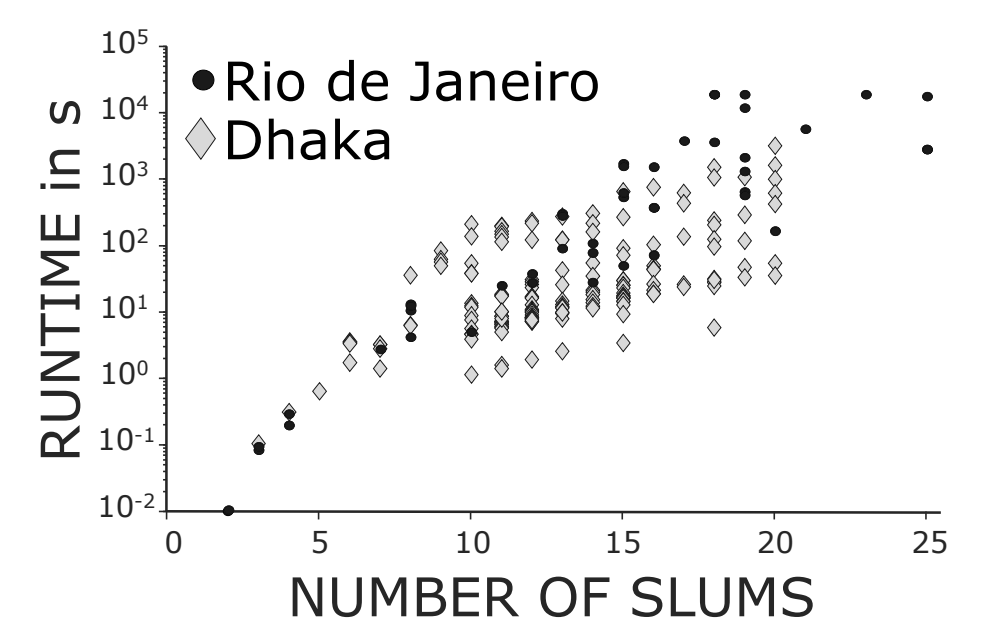
- Clustering of slums in city into subproblems
- Heuristic based on minimal spanning tree and greedy algorithm
- Start solution for MILP solver based on minimal spanning tree



Solution of MILP for slum selection



Heuristic solution for Rio de Janeiro (as tree)



Run time analysis for clustering method

Modeled Mixed Integer Linear Problem

Objective	minimize	Truck costs	+	Fix pipe costs	+	Variable pipe costs	+	Tank costs	
		$\sum_{i \in N_w} \sum_{j \in N} \left(\sum_{k \in K_{\text{truck}}} x_{\text{truck}}(i,j,k) \cdot \text{Cost}_{\text{truck}}(i,j,k) \right)$		$\sum_{k \in K_{\text{pipe}}} (x_{\text{pipe}}(i,j,k) \cdot \text{Cost}_{\text{pipe,fix}}(i,j,k) + x_{\text{exp}}^Q(i,j) \cdot \text{Cost}_{\text{pipe,var}}(i,j,k))$		$\sum_{i \in N} \sum_{k \in K_{\text{tank}}} x_{\text{tank}}(i,k) \cdot \text{Cost}_{\text{tank}}(k)$			
Constraints	Flow condition	s.t. $\sum_{j \in N_w} x^Q(j,i) = \sum_{j \in N} x^Q(i,j) + Q_{\text{daily}}(i)$		$\forall i \in N_w$	Volume flow linearization	$x_p^Q(i,j) = Q_{\text{lin}}(1) + \sum_{m=2}^{ N_{Q^3} } (Q_{\text{lin}}(m) - Q_{\text{lin}}(m-1)) \lambda_{Q_{\text{lin}}}(i,j,k,m)$		$\forall i \in N_w, j \in N, k \in K_{\text{pipe}}$	
	Capacity	$x_p^Q(i,j) + x_t^Q(i,j) = x^Q(i,j)$		$\forall i \in N_w, j \in N$		$x_{\text{exp}}^Q(i,j) = Q_{\text{lin}}^3(1) + \sum_{m=2}^{ N_{Q^3} } (Q_{\text{lin}}^3(m) - Q_{\text{lin}}^3(m-1)) \lambda_{Q_{\text{lin}}}(i,j,k,m)$		$\lambda_{Q_{\text{lin}}}(i,j,k,m) \leq x_{\text{pipe}}(i,j,k)$	$\forall i \in N_w, j \in N, k \in K_{\text{pipe}}, m \in N_{Q^3}$
		$\sum_{k \in K_{\text{truck}}} x_{\text{truck}}(i,j,k) \cdot \text{Capa}_{\text{truck}}(i,j,k) \geq x_t^Q(i,j)$				$\sum_{k \in K_{\text{pipe}}} x_{\text{pipe}}(i,j,k) \cdot \text{Capa}_{\text{pipe}}(i,j,k) \geq x_p^Q(i,j)$			
	Selection relation	$\sum_{j \in N} x^Q(w,j) \leq \text{Capa}_w(w)$		$\forall w \in W$		$\sum_{k \in K_{\text{tank}}} x_{\text{tank}}(i,k) \cdot \text{Capa}_{\text{tank}}(k) \geq x_{\text{Tout}}^Q(i) + x_{\text{TinPout}}^Q(i) + x_{\text{Tindaily}}^Q(i)$		$x_{\text{PInPout}}^Q(i) \geq \sum_{j_1 \in N} x_p^Q(j_1,i) - \sum_{j_2 \in N} x_p^Q(i,j_2)$	
$\sum_{k \in K_{\text{truck}}} x_{\text{truck}}(i,j,k) + \sum_{k \in K_{\text{pipe}}} x_{\text{pipe}}(i,j,k) \geq x_{\text{use}}(i,j)$		$\forall i \in N_w, j \in N$	$\sum_{k \in K_{\text{pipe}}} x_{\text{pipe}}(i,j,k) \leq x_{\text{use}}(i,j)$		$x_{\text{PInPout}}^Q(i) \leq \sum_{j_1 \in N} x_p^Q(j_1,i) - \sum_{j_2 \in N} x_p^Q(i,j_2) + \mathcal{M} \cdot \lambda_{\text{PInPout}}(i)$				
No loops	$\sum_{k \in K_{\text{truck}}} x_{\text{truck}}(i,j,k) \leq x_{\text{use}}(i,j) \cdot N_{\text{truck}}^{\text{max}}$		$\forall i \in N_w$	$x_{\text{pipe}}(i,i) = 0$		$x_{\text{Tindaily}}^Q(i) \geq Q_{\text{daily}}(i) - x_{\text{PInPout}}^Q(i)$		$\forall i \in N$	
	$x_{\text{use}}(i,i) = 0$			$x_{\text{truck}}(i,i) = 0$		$x_{\text{TinPout}}^Q(i) \geq \sum_{j_1 \in N} x_p^Q(i,j_1) - \sum_{j_2 \in N} x_p^Q(j_2,i)$			$\forall i \in N$
	$x_{\text{truck}}(i,i) = 0$				$x_{\text{Tout}}^Q(i) \geq \sum_{j \in N_w} x_t^Q(i,j)$				